Today.





Secret Sharing.



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Correcting for loss or even corruption.

Share secret among *n* people.

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The idea of the day.

Two points make a line. Lots of lines go through one point.

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Polynomials over reals: $a_1, \ldots, a_d \in \Re$, use $x \in \Re$.

Polynomials P(x) with arithmetic modulo p: ¹ $a_i \in \{0, ..., p-1\}$ and

$$P(x) = a_d x^d + a_{d-1} x^{d-1} \cdots + a_0 \pmod{p},$$
 for $x \in \{0, \dots, p-1\}.$

Line: $P(x) = a_1 x + a_0$

Line: $P(x) = a_1x + a_0 = mx + b$





















 $\implies 2x \equiv 1 \pmod{5}$



 $x + 2 \equiv 3x + 1 \pmod{5}$ $\implies 2x \equiv 1 \pmod{5} \implies x \equiv 3 \pmod{5}$ 3 is multiplicative inverse of 2 modulo 5.



 $\implies 2x \equiv 1 \pmod{5} \implies x \equiv 3 \pmod{5}$ 3 is multiplicative inverse of 2 modulo 5. Good when modulus is prime!! Two points make a line.

Fact: Exactly 1 degree $\leq d$ polynomial contains d + 1 points.²

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Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime p contains d + 1 pts.

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Poll.

Two points determine a line. What facts below tell you this?

Say points are $(x_1, y_1), (x_2, y_2)$.

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(A) Line is y = mx + b.

(B) Plug in a point gives an equation: $y_1 = mx_1 + b$

(C) The unknowns are m and b.

(D) If equations have unique solution, done.

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All true.



Why solution? Why unique?

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- (A) Solution cuz: $m = (y_2 y_1)/(x_2 x_1), b = y_1 m(x_1)$
- (B) Unique cuz, only one line goes through two points.
- (C) Try: $(m'x+b') (mx+b) = (m'-m)x + (b-b') = ax + c \neq 0$.
- (D) Either $ax_1 + c \neq 0$ or $ax_2 + c \neq 0$.
- (E) Contradiction.

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Flow poll. (All true. (B) is not a proof, it is restatement.)

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(A) and (D)



















³Points with different x values.





















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Poll:example.

The polynomial from the scheme: $P(x) = 2x^2 + 1x + 3 \pmod{5}$. What is true for the secret sharing scheme using P(x)?

(A) The secret is "2". (B) The secret is "3". (C) A share could be (1,5) cuz P(1) = 5(D) A share could be (2,4)(E) A share could be (0,3)

For a line, $a_1x + a_0 = mx + b$ contains points (1,3) and (2,4).

P(1) =

$$P(1) = m(1) + b \equiv m + b$$

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Backsolve: $b \equiv 2 \pmod{5}$. Secret is 2. And the line is...

 $x+2 \mod 5$.

For a quadratic polynomial, $a_2x^2 + a_1x + a_0$ hits (1,2); (2,4); (3,0).

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$a_2 + a_1 + a_0$	\equiv	2	(mod 5)
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Subtracting 2nd from 3rd yields: $a_1 = 1$. $a_0 = (2 - 4(a_1))2^{-1} = (-2)(2^{-1}) = (3)(3) = 9 \equiv 4 \pmod{5}$

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Quadratic

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So polynomial is $2x^2 + 1x + 4 \pmod{5}$

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Will this always work?

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Multiplicative inverses due to gcd(x,p) = 1, forall $x \in \{1, ..., p-1\}$

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$$P(x) = y_1 \Delta_1(x) + y_2 \Delta_2(x) + \cdots + y_{d+1} \Delta_{d+1}(x).$$

hits points $(x_1, y_1); (x_2, y_2) \cdots (x_{d+1}, y_{d+1})$. Degree *d* polynomial!

Modular Arithmetic Fact: Exactly 1 degree $\leq d$ polynomial with arithmetic modulo prime *p* contains d + 1 pts.

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hits points (x_1, y_1) ; $(x_2, y_2) \cdots (x_{d+1}, y_{d+1})$. Degree *d* polynomial! Construction proves the existence of a polynomial!



Mark what's true.

Poll

Mark what's true.

(A)
$$\Delta_1(x_1) = y_1$$

(B) $\Delta_1(x_1) = 1$
(C) $\Delta_1(x_2) = 0$
(D) $\Delta_1(x_3) = 1$
(E) $\Delta_2(x_2) = 1$
(F) $\Delta_2(x_1) = 0$

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(B), (C), and (E)

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Put the delta functions together.

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Construction proves the existence of the polynomial!

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Proof:

Assume two different polynomials Q(x) and P(x) hit the points.

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Assume two different polynomials Q(x) and P(x) hit the points.

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Must prove Roots fact.





$$\begin{array}{r} 4 \quad x + 4 \\ x - 3 \quad) \quad 4x^2 - 3 \quad x + 2 \\ 4x^2 - 2x \\ ----- \\ 4x + 2 \end{array}$$

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$$4 x + 4 r 4$$

$$x - 3) 4x^{2} - 3 x + 2$$

$$4x^{2} - 2x$$

$$-----$$

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$$-----$$

$$4$$

Divide $4x^2 - 3x + 2$ by (x - 3) modulo 5.

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 $4x^2 - 3x + 2 \equiv (x - 3)(4x + 4) + 4 \pmod{5}$

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 $4x^2 - 3x + 2 \equiv (x - 3)(4x + 4) + 4 \pmod{5}$ In general, divide P(x) by (x - a) gives Q(x) and remainder r.

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In general, divide $P(x)$ by $(x - a)$ gives $Q(x)$ and remainder r .
That is, $P(x) = (x - a)Q(x) + r$
Lemma 1: P(x) has root *a* iff P(x)/(x-a) has remainder 0: P(x) = (x-a)Q(x).

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Lemma 2: P(x) has *d* roots; r_1, \ldots, r_d then $P(x) = c(x - r_1)(x - r_2) \cdots (x - r_d)$. **Proof Sketch:** By induction.

Lemma 1: P(x) has root *a* iff P(x)/(x-a) has remainder 0: P(x) = (x-a)Q(x).

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Lemma 2: P(x) has *d* roots; r_1, \ldots, r_d then $P(x) = c(x - r_1)(x - r_2) \cdots (x - r_d)$. **Proof Sketch:** By induction.

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Roots fact: Any degree *d* polynomial has at most *d* roots.

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- Arithmetic modulo a prime *m* is a **finite field** denoted by F_m or GF(m).
- Intuitively, a field is a set with operations corresponding to addition, multiplication, and division.

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(Almost) the same as what is missing: one P(i).

Runtime.

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Runtime: polynomial in k, n, and $\log p$.

- 1. Evaluate degree k 1 polynomial *n* times using log *p*-bit numbers.
- 2. Reconstruct secret by solving system of *k* equations using log *p*-bit arithmetic.

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Infinite number for reals, rationals, complex numbers!

Two points make a line.

Two points make a line. Compute solution: *m*,*b*.

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Compute solution: *m*,*b*. Unique:

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Unique:

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Today: d + 1 points make a unique degree d polynomial.

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Cuz:

Solution: lagrange interpolation.

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Apply: P(x), Q(x) degree d.

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P(x) - Q(x) is degree $d \implies d$ roots.

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Secret Sharing:

k points on degree k - 1 polynomial is great!

Can hand out *n* points on polynomial as shares.